

# Patterns

### Analysis of complex interconnected data







# Outline

#### Quick Notes

- Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



# Quick Notes

- Reminder, first assignment due in a week
  - <u>http://www.reirab.com/Teaching/NS20/Assignment\_1.pdf</u>
  - Any questions from the description?
  - Join a Group in Mycourses
  - Submit the assignment in Mycourses
  - For assignments, 2<sup>k</sup>% of the grade will be deducted per k days of delay.
- Use slack for easier communications
  - Let me know if you didn't get an invite
- Anyone new in the class?

#### Deadlines

- assignment 1 due on Sep. 20th
- assignment 2 due on Oct. 4th
- assignment 3 due on Oct. 18th
- project proposal slides due on Oct. 25th
- project proposal due on Nov. 1th
- Reviews (first round) due on Nov. 8th
- project progress report due on Nov. 22nd
- Reviews (second round) due on Nov. 29th
- project final report slides due on Dec. 1st
- project final report due on Dec. 6th
- Reviews (third round) due on Dec. 13th
- $\circ~$  project revised report and rebuttal due on Dec. 20th
- $\circ\;$  note: dates are tentative, please check them for the updated deadlines

# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



# Marginals of Adjacency Matrix

- Marginals of  $A \Rightarrow$  Degrees  $\circ d_i = \Sigma_j A_{ij}$
- Sum(A) =  $\Sigma_i \Sigma_j A_{ij} = \Sigma_i d_i = ?$
- If directed indegree and outdegree  $\circ$   $d^{in}_{i} = \Sigma_{j} A_{ji}$  and  $d^{out}_{i} = \Sigma_{j} A_{ij}$

	0	1	2	3	4	5	6	7	8	9	10	11
0	0	1	1	0	0	0	0	0	0	0	0	1
1	1	0	1	1	0	0	0	0	0	0	0	0
2	1	1	0	0	0	0	0	0	0	0	0	0
3	0	1	0	0	1	1	0	0	0	0	0	0
4	0	0	0	1	0	1	1	0	0	0	0	0
5	0	0	0	1	1	0	0	0	0	0	0	0
6	0	0	0	0	1	0	0	1	1	0	0	0
7	0	0	0	0	0	0	1	0	1	0	0	0
8	0	0	0	0	0	0	1	1	0	0	1	0
9	0	0	0	0	0	0	0	0	0	0	1	1
10	0	0	0	0	0	0	0	0	1	1	0	1
11	1	0	0	0	0	0	0	0	0	1	1	0

<u>کو</u>

° (\_\_\_\_\_

# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



# Marginals of Adjacency Matrix

- Marginals of  $A \Rightarrow$  Degrees  $\circ$  d<sub>i</sub> =  $\Sigma_i A_{ii}$
- Sum(A) =  $\Sigma_i \Sigma_i A_{ii} = \Sigma_i d_i = 2E$ If undirected

- mean degree:  $1/n \Sigma_i \Sigma_i A_{ii} = 1/n \Sigma_i d_i$ lacksquare
- Density: **Σ<sub>i</sub>Σ<sub>i</sub>A<sub>ii</sub> / n(n-1)** ullet

	0	1	2	3	4	5	6	7	8	9	10	11
0	0	1	1	0	0	0	0	0	0	0	0	1 ]
1	1	0	1	1	0	0	0	0	0	0	0	0
2	1	1	0	0	0	0	0	0	0	0	0	0
3	0	1	0	0	1	1	0	0	0	0	0	0
4	0	0	0	1	0	1	1	0	0	0	0	0
5	0	0	0	1	1	0	0	0	0	0	0	0
6	0	0	0	0	1	0	0	1	1	0	0	0
7	0	0	0	0	0	0	1	0	1	0	0	0
8	0	0	0	0	0	0	1	1	0	0	1	0
9	0	0	0	0	0	0	0	0	0	0	1	1
10	0	0	0	0	0	0	0	0	1	1	0	1
11	1	0	0	0	0	0	0	0	0	1	1	0

#### Real-world networks are **sparse**

mean degree << N-1 (or E << E<sub>max</sub>)

WWW (Stanford-Berkeley): Social networks (LinkedIn): Communication (MSNIM): Co-authorships (DBLP): Internet (AS-Skitter): Roads (California): Proteins (S.Cerevisiae):

N=319,717 N=6,946,668 N=242,720,596 N=317,080 N=1,719,037 N=1,957,027 N=1,870 (Source: Les

mean degree=9.65 mean degree=8.87 ),596 mean degree=11.1 mean degree=6.62 mean degree=14.91 mean degree=2.82 mean degree=2.39 Leskovec et al., Internet Mathematics, 2009)

From Leskovec's slides

#### Adjacency matrix is filled with zeros!

(Density of the matrix: WWW=1.51\*10<sup>-5</sup>, MSN IM=  $2.27*10^{-8}$ )

#### **Implications?**

° (\* 1997)

#### Real-world networks are **sparse**

mean degree << N-1 (or E << E<sub>mox</sub>)

WWW (Stanford-Berkeley): Social networks (LinkedIn): Communication (MSNIM): Co-authorships (DBLP): Internet (AS-Skitter): Roads (California): Proteins (S.Cerevisiae):

N=319,717 N=6,946,668 N=242,720,596 N=317,080 N=1,719,037 N=1,957,027 N=1,870

mean degree=9.65 mean degree=8.87 mean degree=11.1 mean degree=6.62 mean degree=14.91 mean degree=2.82 mean degree=2.39 (Source: Leskovec et al., Internet Mathematics. 2009)

From Leskovec's slides

#### Adjacency matrix is filled with zeros!

```
(Density of the matrix: WWW=1.51*10<sup>-5</sup>, MSN IM= 2.27*10^{-8})
```

**Implications?** Use sparse representations, density not so informative

# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



# Marginals of Adjacency Matrix

- Marginals of  $A \Rightarrow$  Degrees
  - $\circ$  d<sub>i</sub> =  $\Sigma_{i} A_{ij}$



#### **Degree Distribution**

- shows how many nodes of degree d are in the graph Ο
- degree sequence of all nodes  $\Rightarrow$  count frequencies Ο

# Heavy Tailed Degree Distribution

Degree distribution is often **heavy tailed** in real world networks

There are few nodes with very high degree & many with very small degree

This is often referred to as being **scale-free** 

Degree distribution is almost always plotted in log-log scale



# Example

poisson vs powerlaw degree distribution highways vs airways

In air-traffic networks, we have major hubs and many smaller airports. In highway networks, cities are of comparable connections.



کی کی ا

#### The first observations

Nodes: WWW documents Links: URL links

Over 3 billion documents ROBOT: collects all URL's found in a document and follows them recursively



R. Albert, H. Jeong, A-L Barabasi, *Nature*, 401 130 (1999).

**Network Science: Scale-Free Property** 

# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



#### Power law distribution

Provides a good fit to the linear pattern observed in log-log plots for degree distribution



 $\ln p_d = -\alpha \ln d + \beta$  $p_d = C d^{-\alpha}$ 

What is C?

Power law distribution

 $\ln p_d = -\alpha \ln d + c$  $p_d = C d^{-\alpha}$ 

What is C?

$$C = e^{\beta}$$

Provides a good fit to the linear pattern observed in log-log plots for degree distribution



https://en.wikipedia.org/wiki/Power law



17

# Scale free networks

Networks with power-law degree distribution are coined as scale-free

Since power-law is scale invariance:

$$f(d) = p_d = C d^{-\alpha}$$
$$f(\lambda d) = C \lambda^{-\alpha} d^{-\alpha} = \lambda^{-\alpha} f(d)$$

function f is <u>scale invariance</u> iff

 $f(\lambda \; x)$  =  $\lambda^{\alpha} \; f(x)$  for some a and all  $\lambda$ 

#### Fitting a power law

- Use a log-log scale & fit a line
- CDF is preferred which is also powerlaw  $\Rightarrow$  more accurate exponent

• 
$$p(x=d) = C d^{-\alpha} \Rightarrow p(x=d) \sim C d^{1-\alpha}$$

• Use logarithmic binning





<u>ک</u>

کی ا

# Fitting a power law

- Linear Fit in log-log space •
  - Very good R2 and p-value because of Ο log-log scale!
- Log-Likelihood
  - How likely is function f to fit the data? Ο Allows p-value estimation between two alternatives, there is a tool for this:
  - http://tuvalu.santafe.edu/~aaronc/powerlaws/ Still an active research area



From Cosia's slides

# Mean & variance for a power-law

- well-defined mean only if  $\alpha > 2$
- No finite variance if  $\alpha < 3$ 
  - the degree of a randomly chosen node can be Ο significantly different from the mean degree

- Most real world networks are within this range
  - In the examples datasets of Barbasi book, we can see how 0 variance deviates from expected variance of same mean random network with poisson distribution (dashed green line)



# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



### Powerlaws are normal?

- Income follow a Pareto distribution
  - few individuals earned most of the money & majority earned small amounts
  - in the US 1% of the population earns a disproportionate 15% of the total US income
  - 80/20 rule (<u>Pareto principle</u>): a general rule of thumb

e.g. 20 percent of the code has 80 percent of the errors

• Zipf's law

Performers 20 Percent

High Performers 80 Percent

- distribution of words ranked by their frequency in a random text corpus is approximated by a power-law distribution
- the second item occurs approximately 1/2 as often as the first, and the third item 1/3 as often as the first, and so on



Vilfredo Federico Damaso Pareto (1848 - 1923)



George Kingsley Zipf (1902 - 1950)

# What creates a powerlaw?

Preferential Attachment

a.k.a rich get richer, accumulative advantage, Yule process, Matthew effect

#### Albert Barabasi Model (AB)

- Add one node at the time, add m connections per new node
- the probability of forming a connection to an existing node is **proportional to its degree**





# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



# Marginals of Adjacency Matrix

- Marginals of  $A \Rightarrow$  Degrees  $\bullet$ 
  - $\circ$  d<sub>i</sub> =  $\Sigma_j A_{ij}$

	0	1	2	3	4	5	6	7	8	9	10	11
0	0	1	1	0	0	0	0	0	0	0	0	1
1	1	0	1	1	0	0	0	0	0	0	0	0
2	1	1	0	0	0	0	0	0	0	0	0	0
3	0	1	0	0	1	1	0	0	0	0	0	0
4	0	0	0	1	0	1	1	0	0	0	0	0
5	0	0	0	1	1	0	0	0	0	0	0	0
6	0	0	0	0	1	0	0	1	1	0	0	0
7	0	0	0	0	0	0	1	0	1	0	0	0
8	0	0	0	0	0	0	1	1	0	0	1	0
9	0	0	0	0	0	0	0	0	0	0	1	1
10	0	0	0	0	0	0	0	0	1	1	0	1
11	1	0	0	0	0	0	0	0	0	1	1	0



#### Degree Assortativity

Strong correlation between degree of connecting nodes

- For all edges, look at degrees of endpoints
  - Either nodes tend to connect to similar degree nodes or dissimilar



assortative mixing

# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - **powers of A** & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



# Marginals of Adjacency Matrix

- Marginals of  $A \Rightarrow$  Degrees  $\circ d_i = \Sigma_j A_{ij}$
- Sum(A) =  $\Sigma_i \Sigma_j A_{ij} = \Sigma_i d_i = 2E$ If undirected

	0	1	2	3	4	5	6	7	8	9	10	11
0	0	1	1	0	0	0	0	0	0	0	0	1
1	1	0	1	1	0	0	0	0	0	0	0	0
2	1	1	0	0	0	0	0	0	0	0	0	0
3	0	1	0	0	1	1	0	0	0	0	0	0
4	0	0	0	1	0	1	1	0	0	0	0	0
5	0	0	0	1	1	0	0	0	0	0	0	0
6	0	0	0	0	1	0	0	1	1	0	0	0
7	0	0	0	0	0	0	1	0	1	0	0	0
8	0	0	0	0	0	0	1	1	0	0	1	0
9	0	0	0	0	0	0	0	0	0	0	1	1
10	0	0	0	0	0	0	0	0	1	1	0	1
11	1	0	0	0	0	0	0	0	0	1	1	0

- If directed column-wise & row-wise marginals ⇒ indegree and outdegree
  - $\circ \quad d_{\ i}^{in} = \Sigma_j A_{ji} \text{ and } d_{\ i}^{out} = \Sigma_j A_{ij}$
  - sum(A) = E

° (\_\_\_\_\_

- $A^2$ : # of walks with length two
  - $\circ \qquad \mathbf{A^2}_{ij} = \mathbf{\Sigma}_k \mathbf{A}_{ik} \mathbf{A}_{kj}$
  - If undirected:
    - $A^2_{ij}$ : number of common neighbors
    - What is A<sup>2</sup><sub>ii</sub>? number of neighbors = degree
  - What is A<sup>2</sup> in directed graph? number of reciprocal neighbors







network's reciprocity  $\Sigma_i \Sigma_i A_{ii} A_{ii} / \Sigma_i \Sigma_i A_{ii}$ 

• A<sup>2</sup> : # of walks with length **two** 

- A<sup>3</sup> : # of walks of length **three** 
  - Is it same as number of paths? 0



• A<sup>2</sup> : # of walks with length **two** 

- A<sup>3</sup> : # of walks of length three
  - Is it same as number of paths?
    - A walk is a finite or infinite sequence of edges which joins a sequence of vertices
    - A **trail** is a walk in which all edges are distinct.
    - A **path** is a trail in which all vertices are distinct.

https://en.wikipedia.org/wiki/Path (graph theory)#Walk, trail, path



• A<sup>2</sup> : # of walks with length **two** 

- A<sup>3</sup> : # of walks of length three
  - Is it same as number of paths? **No!**





<u>کا</u>

• A<sup>2</sup> : # of walks with length **two** 

- A<sup>3</sup> : # of walks of length three
  - Is it same as number of paths? **No!**
  - What is  $A_{ii}^3$  ?





• A<sup>2</sup> : # of walks with length **two** 

- A<sup>3</sup> : # of walks of length three
  - Is it same as number of paths? **No!**
  - What is  $A_{ii}^3$  ?

Number of Triangles?





### Example

import networkx as nx	[[0 1 1 1 0] [1 0 1 0 1]	2
<pre>G = nx.random_geometric_graph(5, 0.5)</pre>	[1 1 0 1 1] [1 0 1 0 0]	
<pre>A = nx.adjacency_matrix(G).todense()</pre>	[0 1 1 0 0]]	
print A	[[ <b>3</b> 1 2 1 2] [1 <b>3</b> 2 2 1]	
A2 = A*A	[2 2 <b>4</b> 1 1] [1 2 1 <b>2</b> 1]	3
print A2	[2 1 1 1 <b>2</b> ]]	
A3 = A2 * A	[[ <b>4</b> 7 7 5 3] [7 <b>4</b> 7 3 5] [7 7 <b>6</b> 6 6]	
print A3	[5 3 6 <b>2</b> 3] [3 5 6 3 <b>2</b> ]]	4

Local:  $c_i = A_{ii}^3 / d_i (d_i-1)$ 

Total number of triangles?





<u>ک</u>

° (\* 1

Local:  $c_i = A_{ii}^3 / d_i (d_i-1)$ 

Total number of triangles?



**)** 

کی ا

Local:  $c_i = A_{ii}^3 / d_i (d_i - 1)$ 

Total number of triangles? **Tr(A<sup>3</sup>)/6** 

Real networks have a lot of triangles Friends of friends are friends

Can we compute number of triangles more effectively?





Local:  $c_i = A_{ii}^3 / d_i (d_i-1)$ 

Total number of triangles? **Tr(A<sup>3</sup>)/6** 

Real networks have a lot of triangles Friends of friends are friends

Can we compute number of triangles more effectively? from eigenvalues of A as  $\sum_{i} \lambda_{i}^{3}$ Since if  $\lambda$  is eigenvalue of A then  $\lambda^{p}$  is an eigenvalue of  $A^{p}$ 

We can approximate with using only top eigenvalues since this distribution is skewed Many works on approximating number of triangles in large graphs

6



[[47753]

Local:  $c_i = A_{ii}^3 / d_i (d_i - 1)$ 

Total number of triangles? Tr(A<sup>3</sup>)/6

Global: Tr(A<sup>3</sup>)/(Sum(A<sup>2</sup>)-Tr(A<sup>2</sup>))

measures the density of triangles

 $\frac{(\text{number of closed paths of length 2})}{(\text{number of paths of length 2})}$ 







### Transitivity & Assortativity

• High global clustering coefficient or high average local clustering coefficient

• Distribution of local clustering coefficient



network measure	scope	$\operatorname{graph}$	definition	explanation	
degree	L	U	$k_i = \sum_{j=1}^n A_{ij}$	number of edges attached to ver-	
in-degree	L	D	$k_i^{\text{in}} = \sum_{j=1}^n A_{ji}$	number of arcs terminating at vertex $i$	
out-degree	L	D	$k_i^{\text{out}} = \sum_{j=1}^n A_{ij}$	number of arcs originating from vertex $i$	
edge count	G	U	$m = \frac{1}{2} \sum_{ij} A_{ij}$	number of edges in the network	
arc count	G	D	$m = \sum_{ij} A_{ij}$	number of arcs in the network	
mean degree	G	U	$\langle k \rangle = 2m / n = \frac{1}{n} \sum_{i=1}^{n} k_i$	average number of connections	
mean in- or out-degree	G	D	$\left< k^{\rm in} \right> = \left< k^{\rm out} \right> = 2m  / n$	average number of in- or out- connections per vertex	
reciprocity	G	D	$r = \frac{1}{m} \sum_{ij} A_{ij} A_{ji}$	fraction of directed edges that are reciprocated	
reciprocity	L	D	$r_i = \frac{1}{k_i} \sum_j A_{ij} A_{ji}$	fraction of directed edges from $i$ that are reciprocated	
clustering coefficient	G	U	$c = \frac{\sum_{ijk} A_{ij} A_{jk} A_{ki}}{\sum_{ijk} A_{ij} A_{jk}}$	the network's triangle density	
clustering coefficient	L	U	$c_i = \sum_{jk} A_{ij} A_{jk} A_{ki} \left/ \binom{k_i}{2} \right.$	fraction of pairs of neighbors of $i$ that are also connected	
diameter	G	U	$d = \max_{ij} \ell_{ij}$	length of longest geodesic path in an undirected network	
mean geodesic distance	G	U or D	$\ell = \frac{1}{\binom{n}{2}} \sum_{ij} \ell_{ij}$	$average \ length \ of a \ geodesic \ path$	
eccentricity	G	U or D	$\epsilon_i = \max_i \ell_{ij}$	length of longest geodesic path starting from $i$	<u>From Clauset'</u> <u>slides</u>

Α \_



# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



### Shortest Path

https://en.wikipedia.org/wiki/Breadth-first\_search



Longest & average shortest path

**1** 

کی ا

# Small average shortest path

Shortest path distribution is normal with small [shrinking] average in real world You can reach any node in a graph passing through few hubs This is often referred to as small world

Diameter is also small {longest sp}



Letter-passing experiment, In 1967 discovered the Six Degrees of Separation



Four Degrees of Separation You are 4 hops away from anyone in the planet

**Stanley Milgram** (1933 - 1984)

# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



### Connectivity

**Connected** (undirected) graph: any two vertices can be joined by a path

A **disconnected** graph is made up by two or more connected components



# Connectivity: GCC & bridges

Connected (undirected) graph: any two vertices can be joined by a path A disconnected graph is made up by two or more connected components Largest Component is referred to as the **giant connected component (GCC) Bridge** edges are those that if erased, the graph becomes disconnected



# Connectivity in directed graphs

- **Strongly** connected component
  - has a path from each node to every other node and vice versa
    - e.g. A to B path and B to A path
- Weakly connected component
  - it is connected if we disregard the edge directions

How many scc do we have in this example graph? How many wcc do we have in this example graph?



From Barbasi's slides & From newman's book



# Connectivity in directed graphs

- Strongly connected component
  - has a path from each node to every other node and vice versa
    - e.g. A to B path and B to A path
- Weakly connected component
  - it is connected if we disregard the edge directions

How many scc do we have in this example graph? 5 How many wcc do we have in this example graph? 2



From Barbasi's slides & From newman's book



# In/Out components

In-component: nodes that can reach the scc Out-component: nodes that can be reached from the scc



in/out-component of a specific node: set of nodes reachable by directed paths to/from that node





(b)

out-component of node A

(a)

out-component of node B

From newman's book



# How to check connectivity?

Start from one node, traverse the graph and record the nodes you reach. If the size of this reached set of nodes is equal to all the nodes in the graph, then the graph is connected. If not, this is one component and continue until all nodes have been reached to get all the components.



Depth-first search

Breadth-first search

# Connectivity & Adjacency Matrix

The adjacency matrix of a network with several components can be written in a **block**-diagonal form, so that **nonzero elements are confined to squares**, with all other elements being zero:



From Barbasi's slides

How can we use this to see if the graph is connected based on A?

# Connectivity & Laplacian Matrix

we need to consider a super useful matrix

comes into play in many many different contexts

**D**: diagonal matrix of degrees

### Laplacian Matrix: L = D - A

D	Α
[0 0 0 0 2]]	[0 1 1 0 0]]
1000201	10100
[0 0 4 0 0]	[1 1 0 1 1]
[0 3 0 0 0]	[1 0 1 0 1]
[[3 0 0 0 0]	[[0 1 1 1 0]
	[[3 0 0 0 0] [0 3 0 0 0] [0 0 4 0 0] [0 0 0 2 0] [0 0 0 0 2]]

# Eigenvalues of Laplacian Matrix

Lυ = λυ

wiki: Eigenvalues

L = D - A

[[3-1-1-1 0]

[-1 3 -1 0 -1]

[-1 -1 4 -1 -1] [-1 0 -1 2 0] [0 -1 -1 0 2]]

- We have n eigenvalues which we call Laplacian Spectrum:  $0 = \lambda_0 \le \lambda_1 \le \lambda_2 \le \dots \le \lambda_n$
- $\lambda_0$  is always zero since we have L(1,1...1) = 0 why this holds?
- $E = \frac{1}{2} \Sigma d_{i} = \frac{1}{2} Tr(L) = \frac{1}{2} \Sigma \lambda_{i}$
- Laplacian Spectrum relates to graph connectivity & clustering

for undirected graphs, we can always do eigenvalue decomposition since L is symmetric

## Connectivity & Laplacian Matrix

- smallest eigenvalue of **L** is always zero
- second-smallest eigenvalue of L is called Algebraic connectivity or Fiedler value and is nonzero only if graph is connected
- number of zero eigenvalues of L gives the number of connected components

# Clustering & Laplacian Matrix

- Signs of values in Fiedler eigenvector (associated to Fiedler eigenvalue) tell us how to partition the graph into two components by breaking least edges, i.e. minimum cut solution more on this and spectral clustering later
- eigengap is the difference between subsequent eigenvalues
  - first large eigengap is related to the number of clusters in data
  - first eigengap (=smallest nonzero eigenvalue) is called spectral gap which relates to how quickly the diffusion takes place on the network and density of the graph

See this: <u>https://towardsdatascience.com/spectral-clustering-aba2640c0d5b</u>

# Outline

- Quick Notes
  - Assignment 1, slack
- Adjacency matrix and degree
- Sparsity Pattern
- Scale Free Pattern
  - Power-law degree distribution
  - Fitting a power-law
  - Preferential attachment and AB model
- Assortativity Pattern
- Transitivity Pattern
  - powers of A & counting triangles
- Small world Pattern
  - Shortest path
- Connectivity & eigenvalues of Laplacian matrix
- How to pattern?



### Pattern Detection

- WHY?
  - Understand the language of complex systems
  - Characterize different types of networks
  - Design {efficient} data structure & algorithms
  - Tangled with Measurements, Anomaly detection, Modelling
- HOW?
  - What do networks have in common?
  - How to measure or characterize (nodes, communities, whole) networks?
  - What are universal patterns observed in real world networks?
  - What is structure of real-world networks?

- 1. <u>Newman's collection</u>
- 2. <u>Stanford Large Network</u> <u>Dataset Collection</u>
- 3. <u>The Colorado Index of</u> <u>Complex Networks (ICON)</u>
- 4. <u>The Koblenz Network</u> <u>Collection</u>



From Clauset's slides

کی کی ا

- 1. <u>Newman's collection</u>
- 2. <u>Stanford Large Network</u> <u>Dataset Collection</u>
- 3. <u>The Colorado Index of</u> <u>Complex Networks (ICON)</u>
- 4. <u>The Koblenz Network</u> <u>Collection</u>

Entries found: 668 Networks found: 5333



**1** 

کی ا

- 1. <u>Newman's collection</u>
- 2. <u>Stanford Large Network</u> <u>Dataset Collection</u>
- 3. <u>The Colorado Index of</u> <u>Complex Networks (ICON)</u>
- 4. <u>The Koblenz Network</u> <u>Collection</u>

Let slack know if you come across other large repos

KONECT currently holds 261 networks, of which

- 63 are undirected,
- 107 are directed,
- 91 are bipartite,
- 125 are unweighted,
- 90 allow multiple edges,

- 6 have signed edges,
- 10 have ratings as edges,
- 3 allow multiple weighted edges,
- 18 allow positive weighted edges,
- and 89 have edge arrival times.

- 1. Newman's collection
- 2 Stanford Large Network **Dataset Collection**
- 3. The Colorado Index of Complex Networks (ICON)
- 4. The Koblenz Network Collection

	Affilia	tion	à ctor movies	8-		American Revolution	
ONEC	B-		Club membership	B-		Corporate Leadership	
UNEC	B=		Countries	B=	12	Discons	
	B-		Flickr	B-	-	Livelournal	
	B-		Occupation	B-		Orkut	
- 62	B-		Prosper.com	B-		Record labels	Idoc
• 05	B-		South African Companies	B-		Teams	iges,
• 10	в-		YouTube				as edges.
01							as cages,
• 91		11	Ricon	n#		Cattle	e weighted edges,
1.7			Dolphins	0-		Hans	in tunighted edges
• 12	112		Kandaroo	0.5		Magagues	e weighted edges,
- 00	D+		Dhesus	D+	323	Sheen	a arrival times
• 90	0=		Zehra	24	ste	Succh	e arrivar unles.
		_	Leon	_	_		
	Autho	rship					_
	B-		arXiv cond-mat	B-		DBLP	
	B=	10	Github	B=	-	Producers	
	B=	C	Wikibooks (en)	B=	C	Wikibooks (fr)	
	B=	C	Wikinews (en)	B=	C	Wikinews (fr)	
	B=	C	Wikipedia (de)	B=	C	Wikipedia (en)	
	B=	C	Wikipedia (es)	B=	C	Wikipedia (fr)	
	B	C	Wikipedia (it)	B=	C	Wikiquote (en)	
	BE	0	Wiktionary (de)	BE	U	Wiktionary (en)	
	B=	U	Wiktionary (fr)	в=	_	Writers	
	Citatio	on				- 1.5 C	_
	0-0		arXiv hep-ph	0-0	100	arXiv hep-th	
	0-0		CiteSeer	0-	140	Cora citation	
	D-9	_	DBLP	D= 22		US patents	
	Coaut	horship				Same an	
	0-		arXiv astro-ph	0=	C	arXiv hep-ph	
	U=	C	arXiv hep-th	U=	C	DBLP	
	U=		DBLP co-authorship				_
	Comm	unicati	on				
	D=Q	O	Digg	D=Q	C	DNC emails	
	D=Q	O	Enron	D-Q		EU institution	
	D=Q	O	Facebook	D=Q	C	Linux kernel mailing list replies	
	D=Q	C	Manufacturing emails	D=Q	O	Slashdot	
	U-		🕹 U. Rovira i Virgili	D=	C	UC Irvine messages	
	0-0	m	terle le cli a le	0-0	m	terle le cli	



کی ا



**30** '

کی ک

